

# Ph 219c/CS 219c

## Exercises

Due: Wednesday 9 March 2011

### 3.1 What probability distributions are consistent with a mixed state?

A density operator  $\rho$ , expressed in the orthonormal basis  $\{|\alpha_i\rangle\}$  that diagonalizes it, is

$$\rho = \sum_i p_i |\alpha_i\rangle \langle \alpha_i| . \quad (1)$$

We would like to realize this density operator as an ensemble of pure states  $\{|\varphi_\mu\rangle\}$ , where  $|\varphi_\mu\rangle$  is prepared with a specified probability  $q_\mu$ . This preparation is possible if the  $|\varphi_\mu\rangle$ 's can be chosen so that

$$\rho = \sum_\mu q_\mu |\varphi_\mu\rangle \langle \varphi_\mu| . \quad (2)$$

We say that a probability vector  $q$  (a vector whose components are nonnegative real numbers that sum to 1) is *majorized* by a probability vector  $p$  (denoted  $q \prec p$ ), if there exists a *doubly stochastic* matrix  $D$  such that

$$q_\mu = \sum_i D_{\mu i} p_i . \quad (3)$$

A matrix is doubly stochastic if its entries are nonnegative real numbers such that  $\sum_\mu D_{\mu i} = \sum_i D_{\mu i} = 1$ . That the columns sum to one assures that  $D$  maps probability vectors to probability vectors (*i.e.*, is *stochastic*). That the rows sum to one assures that  $D$  maps the uniform distribution to itself. Applied repeatedly,  $D$  takes any input distribution closer and closer to the uniform distribution (unless  $D$  is a permutation, with one nonzero entry in each row and column). Thus we can view majorization as a partial order on probability vectors such that  $q \prec p$  means that  $q$  is more nearly uniform than  $p$  (or equally close to uniform, in the case where  $D$  is a permutation).

Show that normalized pure states  $\{|\varphi_\mu\rangle\}$  exist such that eq. (2) is satisfied if and only if  $q \prec p$ , where  $p$  is the vector of eigenvalues of  $\rho$ .

(Because the Shannon entropy is Schur concave, it follows that  $H(q) \geq H(p) = H(\rho)$ ; that is, for any realization of the density operator  $\rho$  as an ensemble of pure states  $\{|\varphi_\mu\rangle, q_\mu\}$ , the Shannon entropy of the probability vector  $q$  is at least as large as the Von Neumann entropy of the density operator  $\rho$ .)

**Hint:** Recall that, according to the *Hughston-Jozsa-Wootters Theorem*, if eq. (1) and eq. (2) are both satisfied then there is a unitary matrix  $V_{\mu i}$  such that

$$\sqrt{q_\mu} |\varphi_\mu\rangle = \sum_i \sqrt{p_i} V_{\mu i} |\alpha_i\rangle \quad (4)$$

(see Sec. 2.5.5 of the lecture notes). You may also use (but need not prove) *Horn's Lemma*: if  $q \prec p$ , then there exists a unitary (in fact, orthogonal) matrix  $U_{\mu i}$  such that  $q = Dp$  and  $D_{\mu i} = |U_{\mu i}|^2$ .

### 3.2 Positivity of quantum relative entropy

- a) Show that  $\ln x \leq x - 1$  for all positive real  $x$ , with equality iff  $x = 1$ .  
 b) The (classical) relative entropy of a probability distribution  $\{p(x)\}$  relative to  $\{q(x)\}$  is defined as

$$H(p \parallel q) \equiv \sum_x p(x) (\log p(x) - \log q(x)) . \quad (5)$$

Show that

$$H(p \parallel q) \geq 0 , \quad (6)$$

with equality iff the probability distributions are identical. **Hint:** Apply the inequality from (a) to  $\ln(q(x)/p(x))$ .

- c) The quantum relative entropy of the density operator  $\rho$  with respect to  $\sigma$  is defined as

$$H(\rho \parallel \sigma) = \text{tr } \rho (\log \rho - \log \sigma) . \quad (7)$$

Let  $\{p_i\}$  denote the eigenvalues of  $\rho$  and  $\{q_a\}$  denote the eigenvalues of  $\sigma$ . Show that

$$H(\rho \parallel \sigma) = \sum_i p_i \left( \log p_i - \sum_a D_{ia} \log q_a \right) , \quad (8)$$

where  $D_{ia}$  is a doubly stochastic matrix. Express  $D_{ia}$  in terms of the eigenstates of  $\rho$  and  $\sigma$ . (A matrix is doubly stochastic if its entries are nonnegative real numbers, where each row and each column sums to one.)

d) Show that if  $D_{ia}$  is doubly stochastic, then (for each  $i$ )

$$\log \left( \sum_a D_{ia} q_a \right) \geq \sum_a D_{ia} \log q_a , \quad (9)$$

with equality only if  $D_{ia} = 1$  for some  $a$ .

e) Show that

$$H(\rho \parallel \sigma) \geq H(p \parallel r) , \quad (10)$$

where  $r_i = \sum_a D_{ia} q_a$ .

f) Show that  $H(\rho \parallel \sigma) \geq 0$ , with equality iff  $\rho = \sigma$ .

### 3.3 Properties of Von Neumann entropy

a) Use nonnegativity of quantum relative entropy to prove the *subadditivity* of Von Neumann entropy

$$H(\rho_{AB}) \leq H(\rho_A) + H(\rho_B), \quad (11)$$

with equality iff  $\rho_{AB} = \rho_A \otimes \rho_B$ . **Hint:** Consider the relative entropy of  $\rho_{AB}$  and  $\rho_A \otimes \rho_B$ .

b) Use subadditivity to prove the concavity of the Von Neumann entropy:

$$H\left(\sum_x p_x \rho_x\right) \geq \sum_x p_x H(\rho_x) . \quad (12)$$

**Hint:** Consider

$$\rho_{AB} = \sum_x p_x (\rho_x)_A \otimes (|x\rangle\langle x|)_B , \quad (13)$$

where the states  $\{|x\rangle_B\}$  are mutually orthogonal.

c) Use the condition

$$H(\rho_{AB}) = H(\rho_A) + H(\rho_B) \quad \text{iff} \quad \rho_{AB} = \rho_A \otimes \rho_B \quad (14)$$

to show that, if all  $p_x$ 's are nonzero,

$$H\left(\sum_x p_x \rho_x\right) = \sum_x p_x H(\rho_x) \quad (15)$$

iff all the  $\rho_x$ 's are identical.

d) Use subadditivity to prove the triangle inequality:

$$H(\rho_{AB}) \geq |H(\rho_A) - H(\rho_B)| . \quad (16)$$

**Hint:** Construct a “purification” of  $\rho_{AB}$  — introduce a third system  $C$  and consider  $|\Phi\rangle_{ABC}$  such that

$$\text{tr}_C (|\Phi\rangle\langle\Phi|) = \rho_{AB} ; \quad (17)$$

then use the subadditivity relations  $H(\rho_{BC}) \leq H(\rho_B) + H(\rho_C)$  and  $H(\rho_{AC}) \leq H(\rho_A) + H(\rho_C)$ .

### 3.4 Entanglement of typical bipartite pure states

Suppose that a pure state is chosen at random on the bipartite system  $AB$ , where  $d_A/d_B \ll 1$ . Then with high probability the density operator on  $A$  will be very nearly maximally mixed. The purpose of this problem is to derive this property.

To begin with, we will calculate the value of  $\langle \text{tr} \rho_A^2 \rangle$ , where  $\langle \cdot \rangle$  denotes the average over all pure states  $\{|\varphi\rangle\}$  of  $AB$ , and  $\rho_A = \text{tr}_B (|\varphi\rangle\langle\varphi|)$ .

a) It is convenient to evaluate  $\text{tr} \rho_A^2$  using a trick. Imagine introducing a copy  $A'B'$  of the system  $AB$ . Show that

$$\text{tr}_A \rho_A^2 = \text{tr}_{ABA'B'} [(S_{AA'} \otimes I_{BB'}) (|\varphi\rangle\langle\varphi|)_{AB} \otimes |\varphi\rangle\langle\varphi|_{A'B'}] , \quad (18)$$

where  $S_{AA'}$  denotes the swap operator

$$S_{AA'} : |\varphi\rangle_A \otimes |\psi\rangle_{A'} \mapsto |\psi\rangle_A \otimes |\varphi\rangle_{A'} . \quad (19)$$

b) We wish to average the expression found in (a) over all pure states  $|\varphi\rangle$ . Rather than go into the details of how such an average is defined, I will simply assert that

$$\langle |\varphi\rangle\langle\varphi|_A \otimes |\varphi\rangle\langle\varphi|_{A'} \rangle = C \Pi_{AA'} , \quad (20)$$

where  $C$  is a constant and  $\Pi_{AA'}$  denotes the projector onto the subspace of  $AA'$  that is *symmetric* under interchange of  $A$  and  $A'$ . Eq. (20) can be proved using invariance properties of the average and some group representation theory, but I hope you will regard it as obvious. The state being averaged is symmetric, and the average should not distinguish any symmetric state from any other symmetric state. Express the constant  $C$  in terms of the dimension  $d \equiv d_A = d_{A'}$ .

- c) Use the property  $\Pi_{AA'} = \frac{1}{2}(I_{AA'} + S_{AA'})$  to evaluate the expression found in (a). Show that

$$\langle \text{tr } \rho_A^2 \rangle = \frac{d_A + d_B}{d_A d_B + 1}. \quad (21)$$

- d) Now estimate the average  $L^2$  distance of  $\rho_A$  from the maximally mixed density operator  $\frac{1}{d_A}I_A$ , where  $\|M\|_2 = \sqrt{\text{tr} M^\dagger M}$ ; show that

$$\left\langle \left\| \rho_A - \frac{1}{d_A}I_A \right\|_2 \right\rangle \leq \frac{1}{\sqrt{d_B}}. \quad (22)$$

**Hints:** First estimate  $\left\langle \left\| \rho_A - \frac{1}{d_A}I_A \right\|_2^2 \right\rangle$  using eq. (21) and the obvious property  $\langle \rho_A \rangle = \frac{1}{d_A}I_A$ . Then show that for any nonnegative function  $f$ , it follows from the Cauchy-Schwarz inequality that  $\langle \sqrt{f} \rangle \leq \sqrt{\langle f \rangle}$ , and use this property to estimate  $\left\langle \left\| \rho_A - \frac{1}{d_A}I_A \right\|_2 \right\rangle$ .

- e) Finally, estimate the average  $L^1$  distance of  $\rho_A$  from the maximally mixed density operator, where  $\|M\|_1 = \text{tr} \sqrt{M^\dagger M}$ . Use the Cauchy-Schwarz inequality to show that  $\|M\|_1 \leq \sqrt{d} \|M\|_2$ , if  $M$  is a  $d \times d$  matrix, and that therefore

$$\left\langle \left\| \rho_A - \frac{1}{d_A}I_A \right\|_1 \right\rangle \leq \sqrt{\frac{d_A}{d_B}}. \quad (23)$$

It follows from (d) that the average entanglement entropy of  $A$  and  $B$  is close to maximal for  $d_A/d_B \ll 1$ :  $\langle H(A) \rangle \geq \log_2 d_A - d_A/2d_B \ln 2$ , though you are not asked to prove this bound. Thus, if for example  $A$  is 50 qubits and  $B$  is 100 qubits, the typical entropy deviates from maximal by only about  $2^{-50} \approx 10^{-15}$ .