

Ph 219b/CS 219b

Exercises

Due: Thursday 21 November 2019

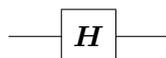
3.1 Universal quantum gates I

In this exercise and the two that follow, we will establish that several simple sets of gates are universal for quantum computation.

The *Hadamard transformation* \mathbf{H} is the single-qubit gate that acts in the standard basis $\{|0\rangle, |1\rangle\}$ as

$$\mathbf{H} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}; \quad (1)$$

in quantum circuit notation, we denote the Hadamard gate as



The single-qubit *phase gate* \mathbf{P} acts in the standard basis as

$$\mathbf{P} = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}, \quad (2)$$

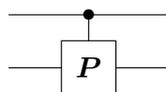
and is denoted



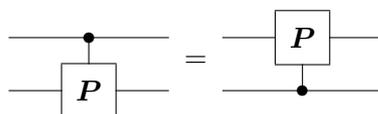
A two-qubit *controlled phase gate* $\Lambda(\mathbf{P})$ acts in the standard basis $\{|00\rangle, |01\rangle, |10\rangle, |11\rangle\}$ as the diagonal 4×4 matrix

$$\Lambda(\mathbf{P}) = \text{diag}(1, 1, 1, i) \quad (3)$$

and can be denoted

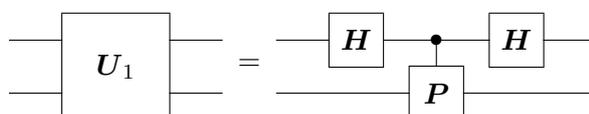


Despite this misleading notation, the gate $\Lambda(\mathbf{P})$ actually acts symmetrically on the two qubits:

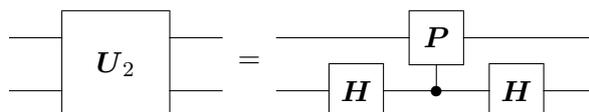


We will see that the two gates \mathbf{H} and $\Lambda(\mathbf{P})$ comprise a *universal gate set* – any unitary transformation can be approximated to arbitrary accuracy by a quantum circuit built out of these gates.

a) Consider the two-qubit unitary transformations \mathbf{U}_1 and \mathbf{U}_2 defined by quantum circuits



and



Let $|ab\rangle$ denote the element of the standard basis where a labels the upper qubit in the circuit diagram and b labels the lower qubit. Write out \mathbf{U}_1 and \mathbf{U}_2 as 4×4 matrices in the standard basis. Show that \mathbf{U}_1 and \mathbf{U}_2 both act trivially on the states

$$|00\rangle, \quad \frac{1}{\sqrt{3}} (|01\rangle + |10\rangle + |11\rangle). \quad (4)$$

b) Thus \mathbf{U}_1 and \mathbf{U}_2 act nontrivially only in the two-dimensional space spanned by

$$\left\{ \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle), \frac{1}{\sqrt{6}} (|01\rangle + |10\rangle - 2|11\rangle) \right\}. \quad (5)$$

Show that, expressed in this basis, they are

$$\mathbf{U}_1 = \frac{1}{4} \begin{pmatrix} 3+i & \sqrt{3}(-1+i) \\ \sqrt{3}(-1+i) & 1+3i \end{pmatrix}, \quad (6)$$

and

$$\mathbf{U}_2 = \frac{1}{4} \begin{pmatrix} 3+i & \sqrt{3}(1-i) \\ \sqrt{3}(1-i) & 1+3i \end{pmatrix}. \quad (7)$$

- c) Now express the action of U_1 and U_2 on this two-dimensional subspace in the form

$$U_1 = \sqrt{i} \left(\frac{1}{\sqrt{2}} - i \frac{1}{\sqrt{2}} \hat{n}_1 \cdot \vec{\sigma} \right), \quad (8)$$

and

$$U_2 = \sqrt{i} \left(\frac{1}{\sqrt{2}} - i \frac{1}{\sqrt{2}} \hat{n}_2 \cdot \vec{\sigma} \right). \quad (9)$$

What are the unit vectors \hat{n}_1 and \hat{n}_2 ?

- d) Consider the transformation $U_2^{-1}U_1$ (Note that U_2^{-1} can also be constructed from the gates \mathbf{H} and $\Lambda(\mathbf{P})$.) Show that it performs a rotation with half-angle $\theta/2$ in the two-dimensional space spanned by the basis eq. (??), where $\cos(\theta/2) = 1/4$.

3.2 Universal quantum gates II

We have now seen how to compose our fundamental quantum gates to perform, in a two-dimensional subspace of the four-dimensional Hilbert space of two qubits, a rotation with $\cos(\theta/2) = 1/4$. In this exercise, we will show that the angle θ is not a rational multiple of π . Equivalently, we will show that

$$e^{i\theta/2} \equiv \cos(\theta/2) + i \sin(\theta/2) = \frac{1}{4} (1 + i\sqrt{15}) \quad (10)$$

is not a root of unity: there is no finite integer power n such that $(e^{i\theta/2})^n = 1$.

Recall that a *polynomial of degree n* is an expression

$$P(x) = \sum_{k=0}^n a_k x^k \quad (11)$$

with $a_n \neq 0$. We say that the polynomial is *rational* if all of the a_k 's are rational numbers, and that it is *monic* if $a_n = 1$. A polynomial is *integral* if all of the a_k 's are integers, and an integral polynomial is *primitive* if the greatest common divisor of $\{a_0, a_1, \dots, a_n\}$ is 1.

- a) Show that the monic rational polynomial of minimal degree that has $e^{i\theta/2}$ as a root is

$$P(x) = x^2 - \frac{1}{2}x + 1. \quad (12)$$

The property that $e^{i\theta/2}$ is not a root of unity follows from the result (a) and the

Theorem *If a is a root of unity, and $P(x)$ is a monic rational polynomial of minimal degree with $P(a) = 0$, then $P(x)$ is integral.*

Since the minimal monic rational polynomial with root $e^{i\theta/2}$ is not integral, we conclude that $e^{i\theta/2}$ is not a root of unity. In the rest of this exercise, we will prove the theorem.

b) By “long division” we can prove that if $A(x)$ and $B(x)$ are rational polynomials, then there exist rational polynomials $Q(x)$ and $R(x)$ such that

$$A(x) = B(x)Q(x) + R(x) , \quad (13)$$

where the “remainder” $R(x)$ has degree less than the degree of $B(x)$. Suppose that $a^n = 1$, and that $P(x)$ is a rational polynomial of minimal degree such that $P(a) = 0$. Show that there is a rational polynomial $Q(x)$ such that

$$x^n - 1 = P(x)Q(x) . \quad (14)$$

c) Show that if $A(x)$ and $B(x)$ are both primitive integral polynomials, then so is their product $C(x) = A(x)B(x)$. **Hint:** If $C(x) = \sum_k c_k x^k$ is not primitive, then there is a prime number p that divides all of the c_k 's. Write $A(x) = \sum_l a_l x^l$, and $B(x) = \sum_m b_m x^m$, let a_r denote the coefficient of lowest order in $A(x)$ that is not divisible by p (which must exist if $A(x)$ is primitive), and let b_s denote the coefficient of lowest order in $B(x)$ that is not divisible by p . Express the product $a_r b_s$ in terms of c_{r+s} and the other a_l 's and b_m 's, and reach a contradiction.

d) Suppose that a monic integral polynomial $P(x)$ can be factored into a product of two monic rational polynomials, $P(x) = A(x)B(x)$. Show that $A(x)$ and $B(x)$ are integral. **Hint:** First note that we may write $A(x) = (1/r) \cdot \tilde{A}(x)$, and $B(x) = (1/s) \cdot \tilde{B}(x)$, where r, s are positive integers, and $\tilde{A}(x)$ and $\tilde{B}(x)$ are primitive integral; then use (c) to show that $r = s = 1$.

e) Combining (b) and (d), prove the theorem.

What have we shown? Since $U_2^{-1}U_1$ is a rotation by an irrational multiple of π , the powers of $U_2^{-1}U_1$ are dense in a $U(1)$ subgroup.

Similar reasoning shows that $\mathbf{U}_1\mathbf{U}_2^{-1}$ is a rotation by the same angle about a different axis, and therefore its powers are dense in another $U(1)$ subgroup. Products of elements of these two noncommuting $U(1)$ subgroups are dense in the $SU(2)$ subgroup that contains both \mathbf{U}_1 and \mathbf{U}_2 .

Furthermore, products of $\Lambda(\mathbf{P})\mathbf{U}_2^{-1}\mathbf{U}_1\Lambda(\mathbf{P})^{-1}$ and $\Lambda(\mathbf{P})\mathbf{U}_1\mathbf{U}_2^{-1}\Lambda(\mathbf{P})^{-1}$ are dense in another $SU(2)$, acting on the span of

$$\left\{ \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle), \frac{1}{\sqrt{6}}(|01\rangle + |10\rangle - 2i|11\rangle) \right\}. \quad (15)$$

Together, these two $SU(2)$ subgroups close on the $SU(3)$ subgroup that acts on the three-dimensional space orthogonal to $|00\rangle$. Conjugating this $SU(3)$ by $\mathbf{H} \otimes \mathbf{H}$ we obtain another $SU(3)$ acting on the three dimensional space orthogonal to $|+, +\rangle$, where $|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$. The only subgroup of $SU(4)$ that contains both of these $SU(3)$ subgroups is $SU(4)$ itself.

Therefore, the circuits constructed from the gate set $\{\mathbf{H}, \Lambda(\mathbf{P})\}$ are dense in $SU(4)$ — we can approximate any two-qubit gate to arbitrary accuracy, which we know suffices for universal quantum computation. Whew!

3.3 Universal quantum gates III

We have shown that the gate set $\{\mathbf{H}, \Lambda(\mathbf{P})\}$ is universal. Thus any gate set from which both \mathbf{H} and $\Lambda(\mathbf{P})$ can be constructed is also universal. In particular, we can see that $\{\mathbf{H}, \mathbf{P}, \Lambda^2(\mathbf{X})\}$ is a universal set.

- a) It is sometimes convenient to characterize a quantum gate by specifying the action of the gate when it conjugates a Pauli operator. Show that \mathbf{H} and \mathbf{P} have the properties

$$\mathbf{H}\mathbf{X}\mathbf{H} = \mathbf{Z}, \quad \mathbf{H}\mathbf{Y}\mathbf{H} = -\mathbf{Y}, \quad \mathbf{H}\mathbf{Z}\mathbf{H} = \mathbf{X}, \quad (16)$$

and

$$\mathbf{P}\mathbf{X}\mathbf{P}^{-1} = \mathbf{Y}, \quad \mathbf{P}\mathbf{Y}\mathbf{P}^{-1} = -\mathbf{X}, \quad \mathbf{P}\mathbf{Z}\mathbf{P}^{-1} = \mathbf{Z}. \quad (17)$$

- b) Note that, since $\mathbf{P}^{-1} = \mathbf{P}^3$, the gate $\mathbf{K} = \mathbf{H}\mathbf{P}^{-1}\mathbf{H}\mathbf{P}\mathbf{H}$ can be constructed using \mathbf{H} and \mathbf{P} . Show that

$$\mathbf{K}\mathbf{X}\mathbf{K} = \mathbf{Y}, \quad \mathbf{K}\mathbf{Y}\mathbf{K} = \mathbf{X}, \quad \mathbf{K}\mathbf{Z}\mathbf{K} = -\mathbf{Z}. \quad (18)$$

c) Construct circuits for $\Lambda^2(\mathbf{Y})$ and $\Lambda^2(\mathbf{Z})$ using the gate set $\{\mathbf{H}, \mathbf{P}, \Lambda^2(\mathbf{X})\}$.

Then complete the proof of universality for this gate set by constructing $\Lambda(\mathbf{P}) \otimes \mathbf{I}$ using $\Lambda^2(\mathbf{X})$, $\Lambda^2(\mathbf{Y})$, and $\Lambda^2(\mathbf{Z})$.

The Toffoli gate $\Lambda^2(\mathbf{X})$ is universal for reversible classical computation. What must be added to realize the full power of quantum computing? We have just seen that the single-qubit gates \mathbf{H} and \mathbf{P} , together with the Toffoli gate, are adequate for reaching any unitary transformation. But in fact, just \mathbf{H} and $\Lambda^2(\mathbf{X})$ suffice to efficiently simulate any quantum computation.

Of course, since \mathbf{H} and $\Lambda^2(\mathbf{X})$ are both real orthogonal matrices, a circuit composed from these gates is necessarily real — there are complex n -qubit unitaries that cannot be constructed with these tools. But a 2^n -dimensional complex vector space is isomorphic to a 2^{n+1} -dimensional real vector space. A complex vector can be encoded by a real vector according to

$$|\psi\rangle = \sum_x \psi_x |x\rangle \mapsto |\tilde{\psi}\rangle = \sum_x (\text{Re } \psi_x) |x, 0\rangle + (\text{Im } \psi_x) |x, 1\rangle, \quad (19)$$

and the action of the unitary transformation \mathbf{U} can be represented by a real orthogonal matrix \tilde{U}_R defined as

$$\begin{aligned} U_R : \quad |x, 0\rangle &\mapsto (\text{Re } U)|x\rangle \otimes |0\rangle + (\text{Im } U)|x\rangle \otimes |1\rangle, \\ |x, 1\rangle &\mapsto -(\text{Im } U)|x\rangle \otimes |0\rangle + (\text{Re } U)|x\rangle \otimes |1\rangle. \end{aligned} \quad (20)$$

To show that the gate set $\{\mathbf{H}, \Lambda^2(\mathbf{X})\}$ is “universal,” it suffices to demonstrate that the real encoding $\Lambda(\mathbf{P})_R$ of $\Lambda(\mathbf{P})$ can be constructed from $\Lambda^2(\mathbf{X})$ and \mathbf{H} .

d) Verify that $\Lambda(\mathbf{P})_R = \Lambda^2(\mathbf{XZ})$.

e) Use $\Lambda^2(\mathbf{X})$ and \mathbf{H} to construct a circuit for $\Lambda^2(\mathbf{XZ})$.

Thus, the classical Toffoli gate does not need much help to unleash the power of quantum computing. In fact, *any* nonclassical single-qubit gate (one that does not preserve the computational basis), combined with the Toffoli gate, is sufficient.

3.4 Universality from any entangling two-qubit gate

We say that a two-qubit unitary quantum gate is *local* if it is a tensor product of single-qubit gates, and that the two-qubit gates \mathbf{U} and \mathbf{V} are *locally equivalent* if one can be transformed to the other by local gates:

$$\mathbf{V} = (\mathbf{A} \otimes \mathbf{B})\mathbf{U}(\mathbf{C} \otimes \mathbf{D}) . \quad (21)$$

It turns out (you are not asked to prove this) that every two-qubit gate is locally equivalent to a gate of the form:

$$\mathbf{V}(\theta_x, \theta_y, \theta_z) = \exp [i (\theta_x \mathbf{X} \otimes \mathbf{X} + \theta_y \mathbf{Y} \otimes \mathbf{Y} + \theta_z \mathbf{Z} \otimes \mathbf{Z})] , \quad (22)$$

where

$$-\pi/4 < \theta_x \leq \theta_y \leq \theta_z \leq \pi/4 . \quad (23)$$

a) Show that $\mathbf{V}(\pi/4, \pi/4, \pi/4)$ is (up to an overall phase) the **SWAP** operation that interchanges the two qubits:

$$\mathbf{SWAP} (|\psi\rangle \otimes |\phi\rangle) = |\phi\rangle \otimes |\psi\rangle . \quad (24)$$

b) Show that $\mathbf{V}(0, 0, \pi/4)$ is locally equivalent to the CNOT gate $\Lambda(\mathbf{X})$.

As discussed in the lecture notes, the CNOT gate $\Lambda(\mathbf{X})$ together with arbitrary single-qubit gates form a universal gate set. But in fact there is nothing special about the the CNOT gate in this regard. *Any* two-qubit gate \mathbf{U} , when combined with arbitrary single-qubit gates, suffices for universality *unless* \mathbf{U} is either local or locally equivalent to **SWAP**.

To demonstrate that \mathbf{U} is universal when assisted by local gates it suffices to construct $\Lambda(\mathbf{X})$ using a circuit containing only local gates and \mathbf{U} gates.

Lemma *If \mathbf{U} is locally equivalent to $\mathbf{V}(\theta_x, \theta_y, \theta_z)$, then $\Lambda(\mathbf{X})$ can be constructed from a circuit using local gates and \mathbf{U} gates except in two cases: (1) $\theta_x = \theta_y = \theta_z = 0$ (\mathbf{U} is local), (2) $\theta_x = \theta_y = \theta_z = \pi/4$ (\mathbf{U} is locally equivalent to **SWAP**).*

You will prove the Lemma in the rest of this exercise.

c) Show that:

$$\begin{aligned}
 (\mathbf{I} \otimes \mathbf{X})\mathbf{V}(\theta_x, \theta_y, \theta_z)(\mathbf{I} \otimes \mathbf{X})\mathbf{V}(\theta_x, \theta_y, \theta_z) &= \mathbf{V}(2\theta_x, 0, 0) , \\
 (\mathbf{I} \otimes \mathbf{Y})\mathbf{V}(\theta_x, \theta_y, \theta_z)(\mathbf{I} \otimes \mathbf{Y})\mathbf{V}(\theta_x, \theta_y, \theta_z) &= \mathbf{V}(0, 2\theta_y, 0) , \\
 (\mathbf{I} \otimes \mathbf{Z})\mathbf{V}(\theta_x, \theta_y, \theta_z)(\mathbf{I} \otimes \mathbf{Z})\mathbf{V}(\theta_x, \theta_y, \theta_z) &= \mathbf{V}(0, 0, 2\theta_z) .
 \end{aligned}
 \tag{25}$$

d) Show that $\mathbf{V}(0, 0, \theta)$ is locally equivalent to the controlled rotation $\Lambda[\mathbf{R}(\hat{n}, 4\theta)]$, where $\mathbf{R}(\hat{n}, 4\theta) = \exp[-2i\theta(\hat{n} \cdot \boldsymbol{\sigma})]$, for an arbitrary axis of rotation \hat{n} . (Here $\boldsymbol{\sigma} = (\mathbf{X}, \mathbf{Y}, \mathbf{Z})$.)

e) Now use the results of (c) and (d) to prove the Lemma.